

26 **2. Main results.** Their paper aims to address the problem, why and when deep
 27 networks can lessen or break the curse of dimensionality. Unlike many related works
 28 that focus on a particular set of functions which have a very special structure (such as
 29 *compositional* or *polynomial*; see Table 1b), they consider functions in the Korobov
 30 spaces which is more general for high dimensional multivariate approximation. The
 31 Korobov spaces are defined by

$$32 \quad (2.1) \quad X^{2,p}(\Omega) = \{f \in L^p(\Omega) : f|_{\partial\Omega} = 0, D^{\mathbf{k}}f \in L^p(\Omega), |\mathbf{k}|_{\infty} \leq 2\},$$

33 with norm $|f|_{2,\infty} = \left\| \frac{\partial^{2d} f}{\partial x_1^2 \dots \partial x_d^2} \right\|_{\infty}$. By establishing a connection with sparse grids,
 34 they present new error estimates for which the curse of dimensionality is lessened; see
 35 Theorem 2.1 below.

36 **THEOREM 2.1.** *For any dimension d and $0 < \varepsilon < 1$, there is a deep ReLU net-*
 37 *work with d inputs x_1, \dots, x_d capable of expressing any function f in $X^{2,p}([0, 1]^d)$*
 38 *that satisfies $|f|_{2,\infty} \leq 1$ with accuracy ε , and has depth $\mathcal{O}(|\log_2 \varepsilon|(d-1))$ and size*
 39 *$\mathcal{O}(\varepsilon^{-\frac{1}{2}} |\log_2 \varepsilon|^{\frac{3}{2}(d-1)+1} (d-1))$.*

40 Compared with the result in Table 1a, the exponent d only affects logarithm
 41 factors $|\log_2 \varepsilon|$ instead of ε^{-1} . Their result shows that Deep ReLU networks can
 42 significantly lessen the effect of large dimensions d .

43 **3. Sketch of proof.** They use the following technique to prove the new error
 44 bounds. Show certain functions f can be approximated by sparse grids f_m to any
 45 prescribed accuracy ε , and so sparse grids f_M by neural networks f_N of size N .
 46 Together, the approximation error can be decomposed as

$$47 \quad (3.1) \quad \|f - f_N\| \leq \|f - f_m\| + \|f_m - f_N\|,$$

48 for some norm $\|\cdot\|$.

49 **3.1. Approximating functions in the Korobov spaces using sparse grids.**

50 To approximate functions of d variables $\mathbf{x} = (x_1, \dots, x_d) \in [0, 1]^d$, they introduce a
 51 tensor product construction. One can consider a family of grids $\Omega_{\mathbf{l}}$ of level $\mathbf{l} =$
 52 (l_1, \dots, l_d) with a grid size $\mathbf{h}_{\mathbf{l}} = (2^{-l_1}, \dots, 2^{-l_d})$ and $2^{\mathbf{l}} - 1$ points $\mathbf{x}_{\mathbf{i},\mathbf{l}} = \mathbf{i} \otimes \mathbf{h}_{\mathbf{l}}$,
 53 $\mathbf{1} < \mathbf{i} < \mathbf{2}^{\mathbf{l}} - \mathbf{1}$. For each $\Omega_{\mathbf{l}}$, one defines piecewise linear hat functions

$$54 \quad (3.2) \quad \phi_{\mathbf{l},\mathbf{i}} = \prod_{j=1}^d \phi_{l_j, i_j}(x_j),$$

55 where $\phi(x_{l,i}) = \phi(\frac{x-x_{l,i}}{h_l})$ and $\phi(x) = \max(0, 1 - |x|)$.

56 Consider a function spaces spanned by these functions $V_l = \text{span}\{\phi_{\mathbf{l},i} : \mathbf{1} \leq$
 57 $\mathbf{i} \leq \mathbf{2}^l - 1\}$ and the hierarchical increments space $W_l = \text{span}\{\phi_{\mathbf{l},i} : \mathbf{i} \in \mathbf{I}_l\}$, where
 58 $\mathbf{I}_l = \{\mathbf{i} \in \mathbb{N}^d : \mathbf{1} \leq \mathbf{i} \leq \mathbf{2}^l - 1, i_j \text{ odd for all } j\}$. These increment spaces satisfy the
 59 relation, $V_{\mathbf{m}} = \bigoplus_{\mathbf{1} \leq \mathbf{l} \leq \mathbf{m}} W_l$.

60 Sparse grids are discretizations of $X^{2,p}(\Omega)$ defined by $V_m^{(1)} = \bigoplus_{\mathbf{1} \leq \mathbf{l}_1 \leq m+d-1} W_l$ and
 61 correspond to a number of grid points $M = \mathcal{O}(h_m^{-1} |\log_2 h_m|^{d-1})$; see Figure 1 for a
 62 sparse grid in two dimensions. For any $f_m^{(1)} \in V_m^{(1)}$,

$$63 \quad (3.3) \quad f_m^{(1)}(\mathbf{x}) = \sum_{|\mathbf{l}_1 \leq m+d-1} \sum_{i \in \mathbf{I}_l} v_{\mathbf{l},i} \phi_{\mathbf{l},i}(\mathbf{x}),$$

64 where the hierarchical coefficients $v_{\mathbf{l},i}$ depends on two order mixed derivatives of f .
 65 For any prescribed accuracy ε , $\|f - f_m^{(1)}\|_{\infty} = \varepsilon$ with $N = \mathcal{O}(\varepsilon^{-\frac{1}{2}} |\log_2 \varepsilon|^{\frac{3}{2}(d-1)})$.

66 **3.2. Approximating sparse grids by deep networks.** The following propo-
 67 sition shows how deep networks can approximate multidimensional hat functions.

68 PROPOSITION 3.1. *For any dimension d and $0 < \varepsilon < 1$, there is a deep ReLU net-*
 69 *work with d inputs x_1, \dots, x_d that estimates the multiplication $\phi_{\mathbf{l},i}(\mathbf{x}) = \prod_{j=1}^d \phi_{l_j, i_j}(x_j)$*
 70 *with accuracy ε , outputs 0 if one of the $\phi_{l_j, i_j}(x_j)$ is 0, and has depth $\mathcal{O}(|\log_2 \varepsilon| |\log_2 d|)$*
 71 *and size $\mathcal{O}(|\log_2 \varepsilon|(d-1))$.*

72 Then, with the fact that functions in $X^{2,p}([0, 1]^d)$ can be approximated by sparse
 73 grids $f_m \in V_m^{(1)}$, show that sparse grids can be represented by deep networks f_N using
 74 the approximated multiplication written as $\tilde{\phi}_{\mathbf{l},i}(\mathbf{x})$:

$$75 \quad (3.4) \quad f_N(\mathbf{x}) = \sum_{|\mathbf{l}_1 \leq m+d-1} \sum_{i \in \mathbf{I}_l} v_{\mathbf{l},i} \tilde{\phi}_{\mathbf{l},i}(\mathbf{x}).$$

76 The corresponding network is shown in Figure 2.

77 **4. Conclusion.** Their proof is based on the ability of deep networks to approx-
 78 imate sparse grids via a binary tree structure (see Figure 2a). Their result provides
 79 an upper bound for the approximation complexity when the same network is used
 80 to approximate all functions in a given Korobov space, without taking advantage of
 81 special properties of the approximated functions. Yet it is pointed out that sparse
 82 grids they used are highly *anisotropic*: to be efficient, these require the functions
 83 being approximated to be aligned with the axes.

Table 1: Approximation results for different activation functions.

(a) Approximation results with the curse of dimensionality.

	Shallow	Deep
$\sigma \in C^\infty(\mathbb{R})$ (not polynomial)	$f \in W^{m,p}([-1, 1]^d)$ depth 1, size $\mathcal{O}(\varepsilon^{-\frac{d}{m}})$ $\ \cdot\ _p$	-
$\sigma \in C^\infty(\mathbb{R})$ (not polynomial)	f analytic in E_ρ depth 1, size $\mathcal{O}(\log_\rho \varepsilon)$ $\ \cdot\ _p$	-
σ ReLU	$f \in W^{m,2}(B^d)$ depth 1, size $\mathcal{O}(\varepsilon^{-\frac{d}{m}})$ $\ \cdot\ _2$	$f \in W^{m,\infty}([0, 1]^d)$ depth $\mathcal{O}(\log_2 \varepsilon)$, size $\mathcal{O}(\varepsilon^{-\frac{d}{m}}) \log_2 \varepsilon $ $\ \cdot\ _\infty$

(b) Approximation results without the curse of dimensionality.

	Shallow	Deep
$\sigma \in C^\infty(\mathbb{R})$ (not polynomial)	$f \in W^{m,\infty}([-1, 1]^d)$, compositional depth 1, size $\mathcal{O}(\varepsilon^{-\frac{d}{m}})$ $\ \cdot\ _\infty$	$f \in W^{m,\infty}([-1, 1]^d)$, compositional depth $\log_2 d$, size $\mathcal{O}((d-1)\varepsilon^{-\frac{d}{m}})$ $\ \cdot\ _\infty$
σ ReLU	f Lipschitz, $[-1, 1]^d$, compositional depth 1, size $\mathcal{O}(\varepsilon^{-d})$ $\ \cdot\ _\infty$	f Lipschitz, $[-1, 1]^d$, compositional depth $\log_2 d$, size $\mathcal{O}((d-1)\varepsilon^{-w})$ $\ \cdot\ _\infty$

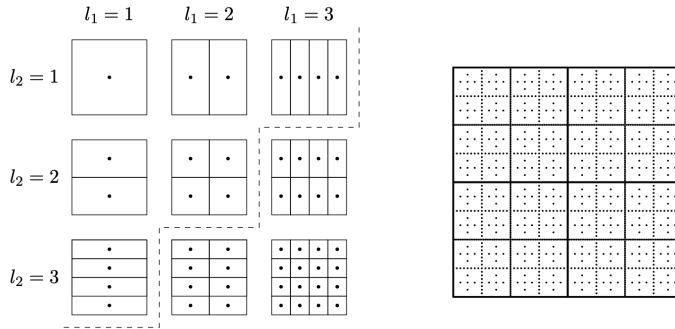
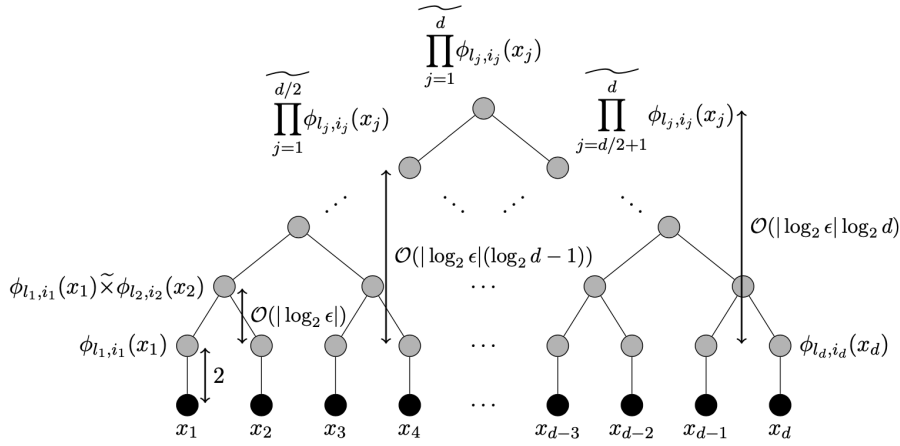
Fig. 1: Left: All subspaces W_l in two dimensions for $(l_1, l_2) \leq (3, 3)$, and sparse and full grids $V_3^{(1)}$ and $V_3^{(\infty)}$. Right: A sparse grid in two dimensions.

Fig. 2: The sparse grid based deep network.

(a) The network that implements the $(d - 1)$ products in $\prod_{j=1}^d \phi_{l_j, i_j}(x_j)$ with a binary tree structure.



(b) The network consists of M subnetworks S_1, S_2, \dots, S_M , which implement the multiplication, $\prod_{j=1}^d \phi_{l_j, i_j}(x_j)$.

